

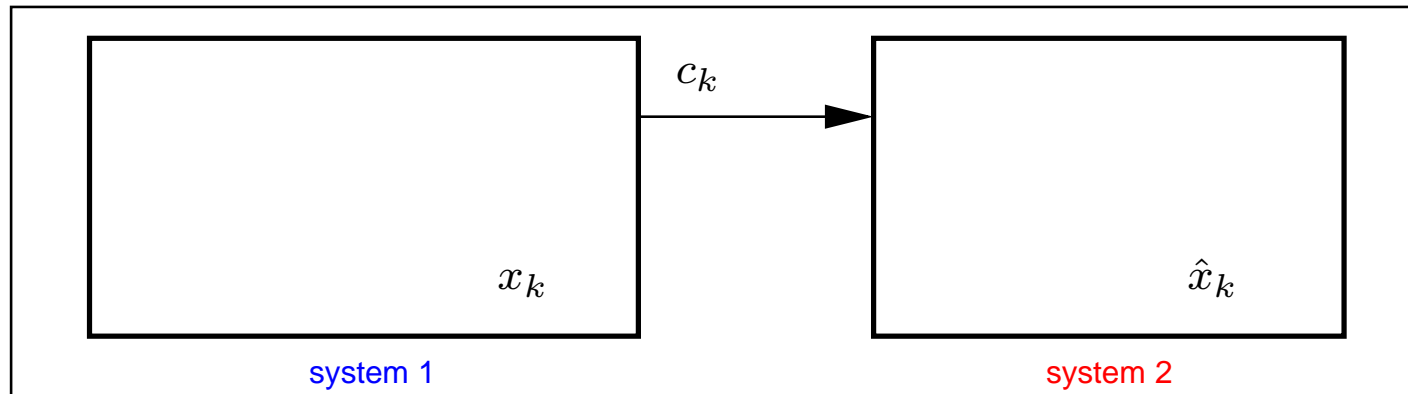
# Synchronization of complex dynamics: when control theory enters the scene of cryptography

Gilles Millérioux

Université Henri Poincaré (Nancy, France)

Centre de Recherche en Automatique de Nancy

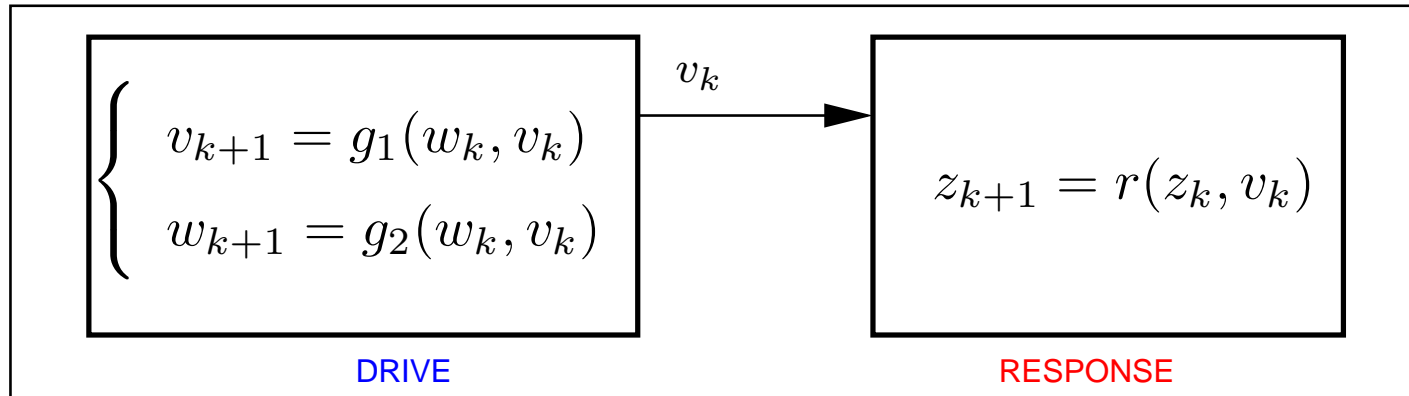
- Synchronization of complex dynamics
- Polytopic formulation
- Control theory and cryptography
- Video



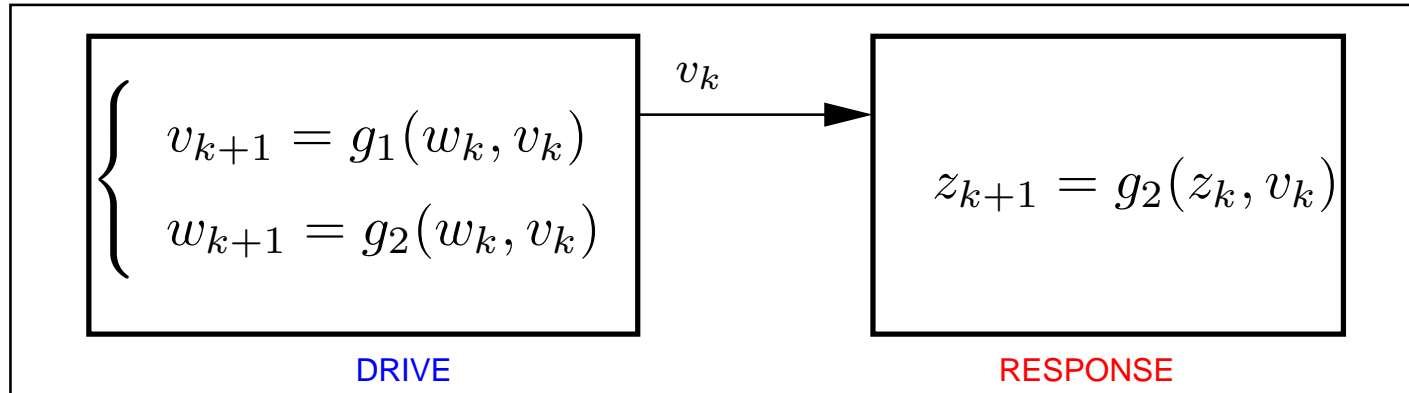
**Definition:** Systems 1 and 2 are self-synchronized if there exists a function  $f$  so that  $f(x_k, \hat{x}_k) = 0$  without the help of any external signal forcing system 2 but the coupling  $c_k$

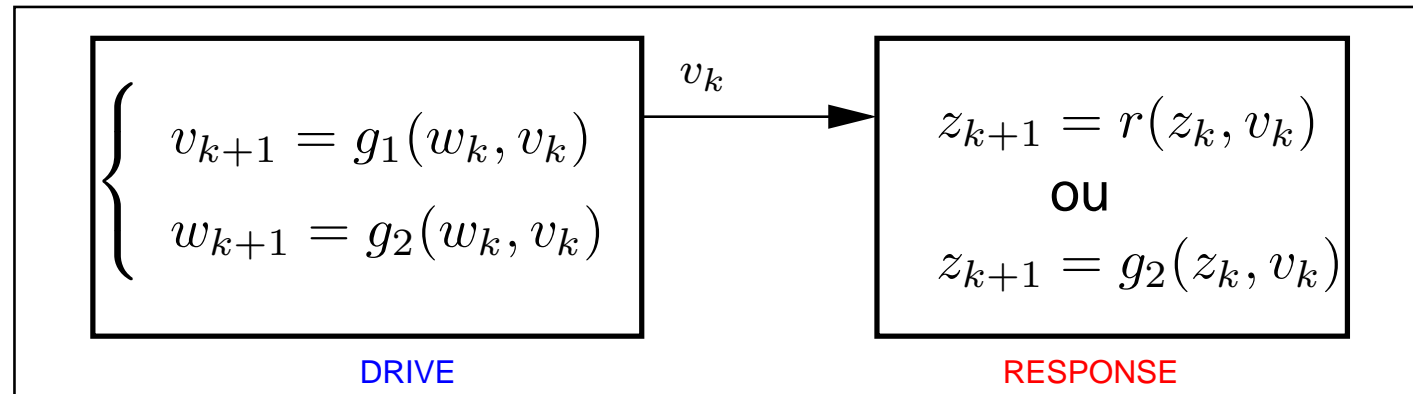
$$\Rightarrow \text{for example } \|x_k - \hat{x}_k\| = 0 \text{ or } \|x_k - \hat{x}_{k+r}\| = 0$$

- heterogeneous configuration ( $r \neq g_2$ )



- homogeneous configuration ( $r = g_2$ )

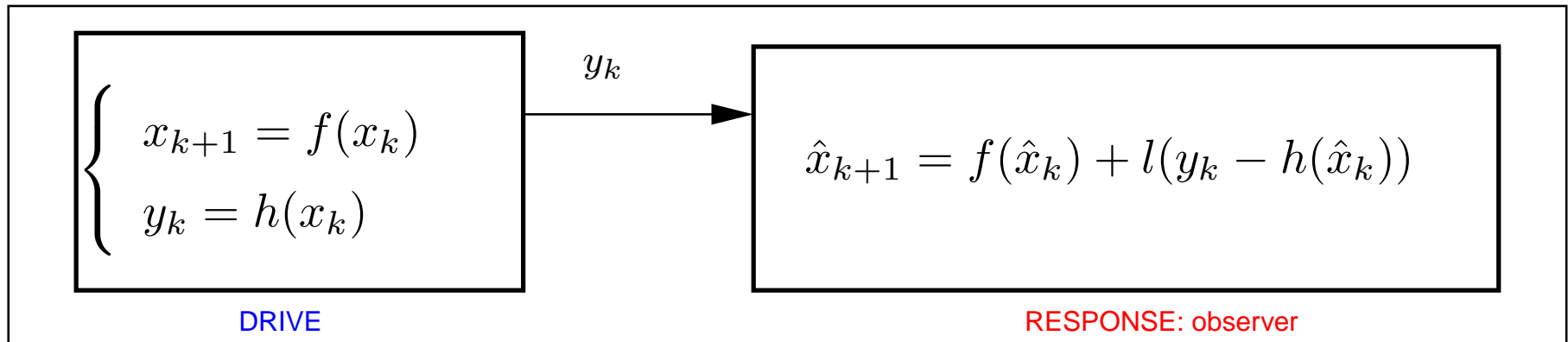




## Chaos synchronization issue (first formulation)

- finding out  $r$  or  $g_{1,2}$  of the DRIVE
- finding out a drive signal  $v_k$

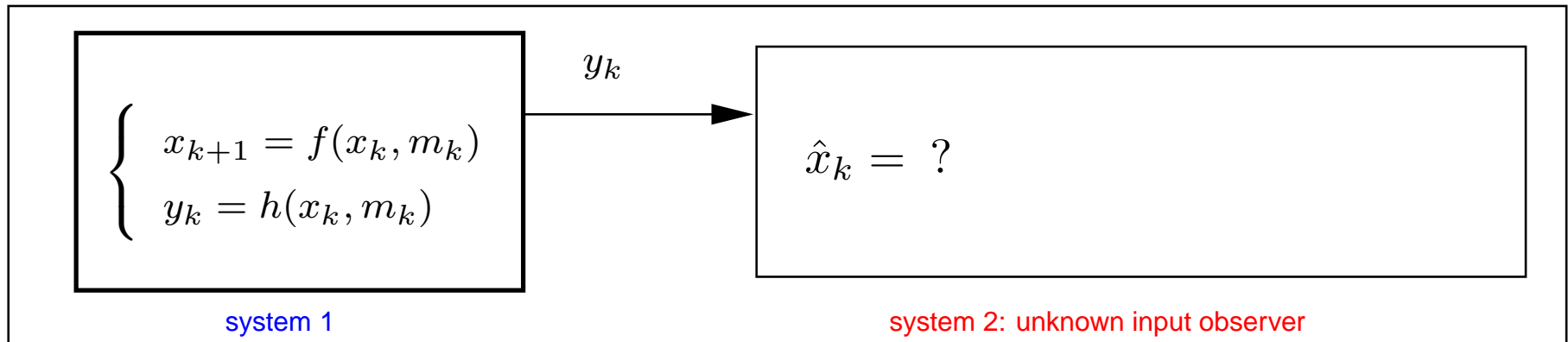
$$\lim_{k \rightarrow \infty} \|z_k - w_k\| = 0$$



## Chaos synchronization issue (observer formulation)

- finding out a function  $l$
- finding out an output  $y_k$

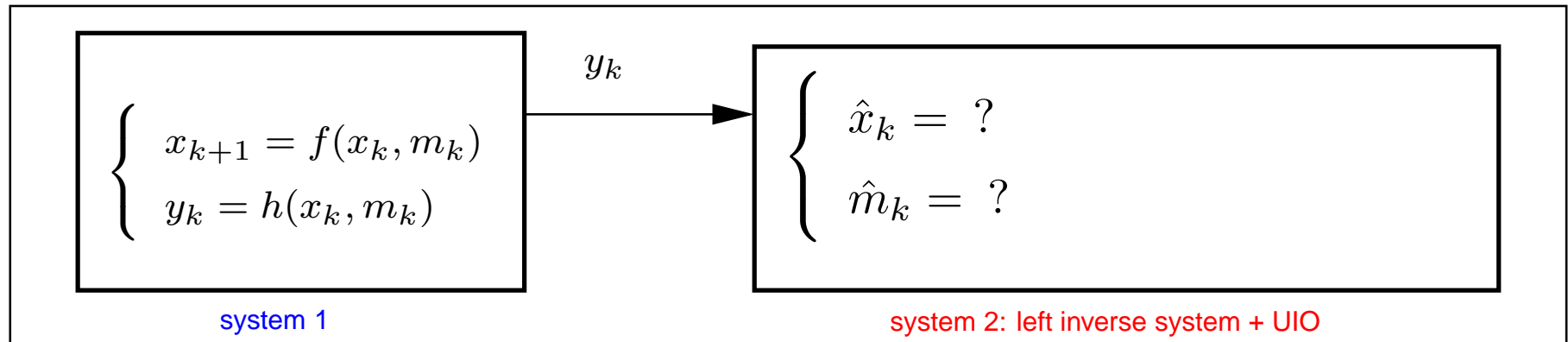
$$\lim_{k \rightarrow \infty} \|\hat{x}_k - x_k\| = 0$$



## Finite time self-synchronization with unknown inputs

- finding out a system 2
- finding out an appropriate output  $y_k$

$$\exists k_f < \infty, \forall \hat{x}_0 \in U, \forall k > k_f \text{ and } \forall m_k \text{ s.t. } \|x_k - \hat{x}_k\| = 0$$



## Finite time self-synchronization with unknown inputs

- finding out a system 2
- finding out an appropriate output  $y_k$

$$\exists k_f < \infty, \forall \hat{x}_0 \in U, \forall k > k_f \text{ and } \forall m_k \text{ s.t.}$$

$$\|x_k - \hat{x}_k\| = 0 \text{ and } \|\hat{m}_k - m_k\| = 0$$



- Switched linear systems (Tent Map, Lozi Map)

$$\begin{cases} x_{k+1} &= A_i x_k \\ y_k &= C x_k \end{cases}$$

- Polynomial systems (Henon Map, Burger Map)

$$\begin{cases} x_{k+1} &= p(x_k, x_k^2, \dots, x_k^{n-1}) x_k \\ y_k &= C x_k \end{cases}$$



LPV system with $\rho_k = h(x_k)$	$\begin{cases} x_{k+1} &= \mathcal{A}(\rho_k) x_k \\ y_k &= C x_k \end{cases}$
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⇒ Polytopic observers with LMI-based design

## Polytopic decomposition principle: example

Drive system

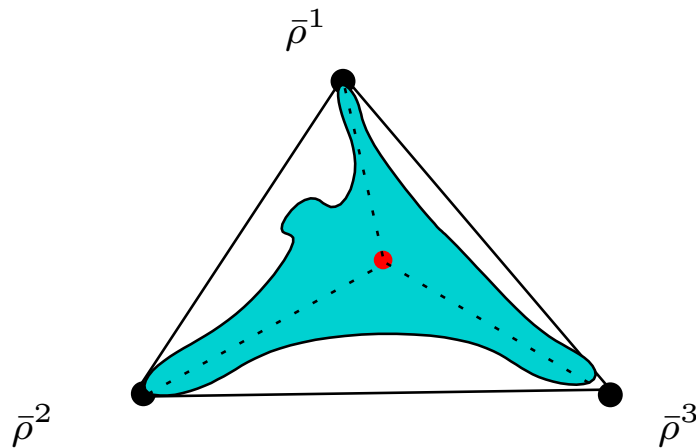
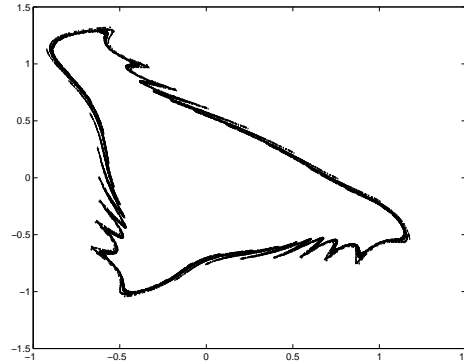
$$\begin{cases} v_{k+1} = -w_k + \alpha(-0.78v_k - 0.28v_k^2 + 0.98v_k^2 \cdot w_k(1 - v_k^2)) \\ w_{k+1} = (1 + 0.12\alpha)w_k \end{cases}$$

$$x_{k+1} = \begin{bmatrix} \alpha(-0.78 - 0.28v_k) & -1 + \alpha 0.98v_k^2 \cdot (1 - v_k^2) \\ 0 & (1 + 0.12\alpha) \end{bmatrix} x_k$$

$$= \begin{bmatrix} \rho_k^1 & \rho_k^2 \\ 0 & (1 + 0.12\alpha) \end{bmatrix} x_k$$

with  $x_k = [v_k \ w_k]^T$

# Polytopic formulation



$$\Rightarrow \rho_k = \sum_{i=1}^3 \xi_k^i(\rho_k) \bar{\rho}^i$$

with  $\sum_{i=1}^3 \xi_k^i = 1$  and  $\xi_k^i \geq 0$

$$\Rightarrow f(x_k) = \mathcal{A}(\rho_k)x_k = \sum_{i=1}^N \xi_k^i A_i x_k$$

Recall of the general description ( $\epsilon_k = x_k - \hat{x}_k$ )

$$\epsilon_{k+1} = f(x_k) - f(\hat{x}_k) - l(y_k - h(\hat{x}_k))$$

Polytopic description (strictly equivalent)

$$\epsilon_{k+1} = \sum_{i=1}^N \xi_k^i A_i x_k - \sum_{i=1}^N \xi_k^i A_i \hat{x}_k - \sum_{i=1}^N \xi_k^i L_i C (x_k - \hat{x}_k)$$

The state reconstruction error obeys

$$\epsilon_{k+1} = \sum_{i=1}^N \xi_k^i (A_i - L_i C) \epsilon_k$$

⇒ stability of LPV polytopic systems

**Theorem 1**  $\epsilon_k$  globally converges to zero if there exist some symmetric matrices  $P_i$ , matrices  $G_i$  and  $F_i$  such that,  $\forall (i, j) \in \{1, \dots, N\} \times \{1, \dots, N\}$ , the following set of Linear Matrix Inequalities is feasible

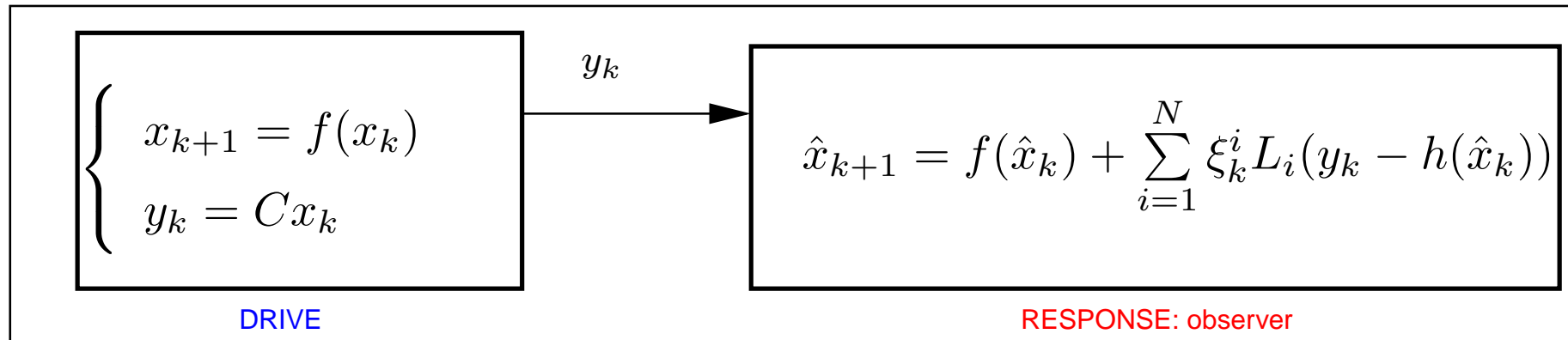
$$\begin{bmatrix} P_i & (G_i A_i - F_i C)^T \\ G_i A_i - F_i C & G_i + G_i^T - P_j \end{bmatrix} > 0 \quad (1)$$

and  $L_i = G_i^{-1} F_i$

The LMIs ensure the existence of a Lyapunov function

$V : \mathbb{R}^n \rightarrow \mathbb{R}^+$  which fulfills :

$$V(\epsilon_{k+1}) - V(\epsilon_k) < 0 \quad \forall k \quad (2)$$

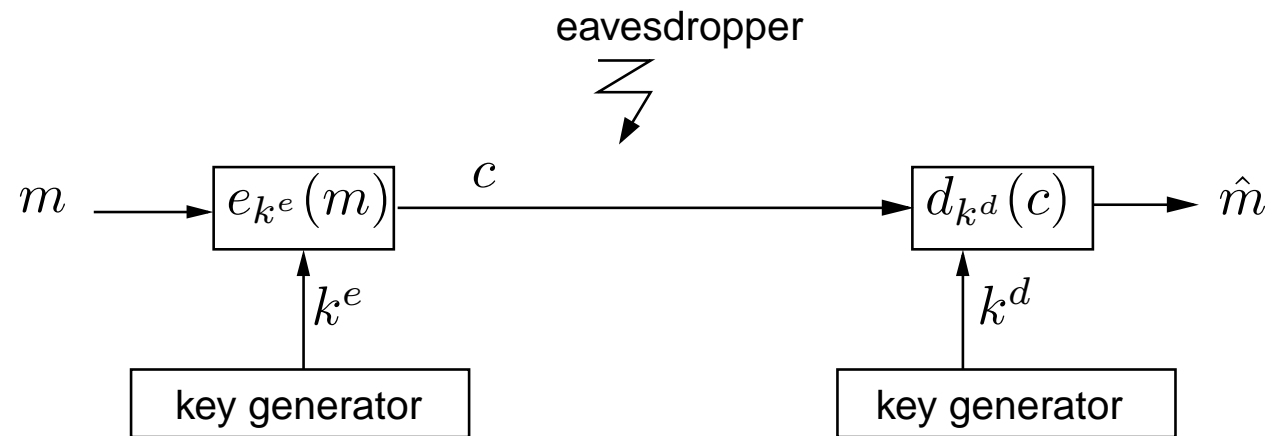


## Chaos synchronization issue (polytopic formulation)

- finding out gains  $L_i$  (solving tractable LMIs)
- finding out a matrix  $C$

$$\lim_{k \rightarrow \infty} \|\hat{x}_k - x_k\| = 0$$

## General ciphering setup



Cryptography must guarantee at least:

- confidentiality

## Classes of ciphers

public-key

private-key

block ciphers

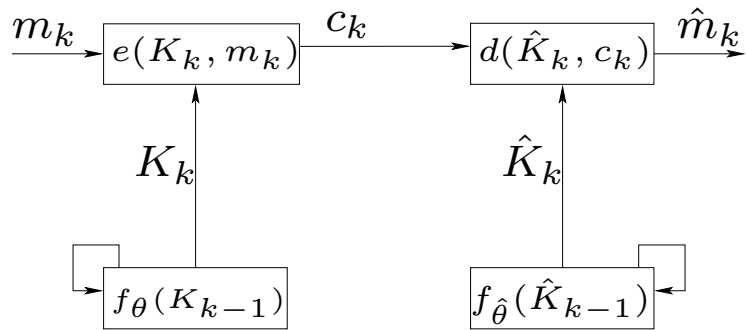
stream ciphers

synchronous

self-synchronous

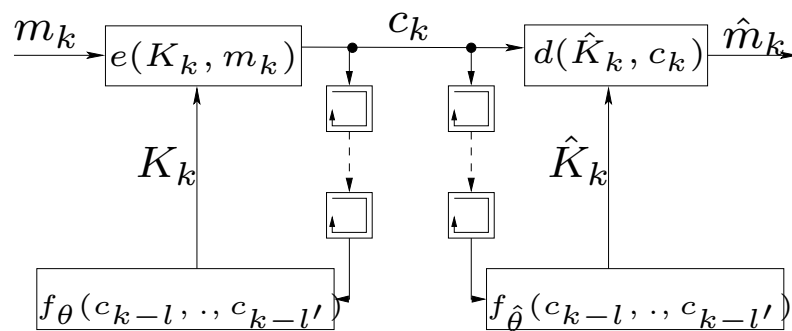


## Stream ciphers



$$\begin{cases} K_k = f_\theta(K_{k-1}) \\ c_k = e(K_k, m_k) \end{cases}$$

Synchronizing Stream Cipher (SSC)



$$\begin{cases} K_k = g_\theta(c_{k-l}, \dots, c_{k-l'}) \\ c_k = e(K_k, m_k) \end{cases}$$

Self Synchronizing Stream Cipher (SSSC)

Relative degree, left invertibility

$$\begin{cases} x_{k+1} = f_{\theta}(x_k, m_k) \\ y_k = h_{\theta}(x_k, [m_k]) \end{cases} \quad (3)$$

The **relative degree** of the system with respect to the quantity  $m_k$  is the required number  $r$  of iterations of the output  $y_k$  so that  $y_{k+r}$  depends on  $m_k$  which actually appears explicitly in the expression of  $y_{k+r}$

$$\begin{aligned} y_{k+r} &= h(f_{\theta}^r(x_k, m_k)) \text{ for } r > 1 \\ &= h(x_k, m_k) \text{ for } r = 0 \end{aligned} \quad (4)$$

The dynamical system is **left invertible** if there exists a nonnegative integer  $R < \infty$ , called *inherent delay*, so that for any two inputs  $m_k$  and  $m'_k$  and any state  $x_k$  the following inference holds:

$$\begin{aligned} h^{(0)}(x_k, m_k) \cdots h^{(R)}(x_k, m_k \cdots m_{k+R}) &= h^{(0)}(x_k, m'_k) \cdots h^{(R)}(x_k, m'_k \cdots m'_{k+R}) \\ \Rightarrow m_k &= m'_k. \end{aligned} \quad (5)$$

## Flatness

$$\begin{cases} x_{k+1} = f_{\theta}(x_k, m_k) \\ y_k = h_{\theta}(x_k, [m_k]) \end{cases} \quad (6)$$

The system with dynamics  $f$ , input  $m_k$  and state vector  $x_k$  of dimension  $n$  is said to be **flat** if there exists an output  $y_k$ , referred to as a flat output, such that all system variables can be expressed as a function of the flat output and a finite number of its backward and/or forward iterates.

$$\begin{cases} x_k = \mathcal{F}_{\theta}(y_{k+k_1}, \dots, y_{k+k'_1}) \\ m_k = \mathcal{G}_{\theta}(y_{k+k_2}, \dots, y_{k+k'_2}) \end{cases} \quad (7)$$

If the DRIVE has  $r < \infty$ ,  $R < \infty$  and is *flat*

**DRIVE  $\Leftrightarrow$  SSSC**

$$\begin{array}{lcl} x_{k+1} & = & f_{\theta}(x_k, m_k) \\ y_k & = & h(x_k, [m_k]) \end{array} \Leftrightarrow \begin{array}{lcl} x_k & = & \mathcal{F}_{\theta}(y_{k+k_1}, \dots, y_{k+k'_1}) \\ y_{k+r} & = & h(f_{\theta}^r(x_k, m_k)) = l(x_k, m_k) \end{array}$$

Identifying with

$$\left\{ \begin{array}{l} K_k = g_{\theta}(c_{k-l}, \dots, c_{k-l'}) \\ c_k = e(K_k, m_k) \end{array} \right. \quad \text{gives}$$

- key generator  $g_{\theta} = \mathcal{F}_{\theta}$
- running key  $K_k = x_k$ ,
- ciphertext  $c_k = y_{k+r}$
- encrypting function  $e = l$

- Observer-based approaches for self-synchronization
- Polytopic formulation defines an appropriate and general framework for chaotic systems
- Dynamical systems make sense in cryptography when defined over a finite field under the flatness condition

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